"Lockless" Get() in RocksDB? Scale with CPU Cores

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100 micros

100 nanos







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Transportation Security Administration















• Mutex becomes extremely expensive when storage access is fast



Under the Mutex

• What is done inside mutex?

mutex_.lock() auto* sv = SuperVersion.ref() mutex_.unlock()

{ data retrievel }

mutex_.lock() if (sv->unref()) { delete sv mutex_.unlock()

// ++int64_t & pointer assignment

Observation

SuperVersion change is a relatively infrequent event

mutex_.lock()
auto* sv = SuperVersion.ref()
mutex_.unlock()

{ data retrievel }

mutex_.lock()
if (sv->unref()) {
 delete sv
}
mutex_.unlock()

Compaction



Proposed Solution

SuperVersion change is a relatively infrequent event

auto* sv = ThreadLocal.Get()

if (UNLIKELY(sv->version != global version)) { mutex .lock() auto* sv = SuperVersion.ref() mutex .unlock()

{ data retrievel }

ThreadLocal.Put(sv)



CPU Breakdown on 32 Core Server



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% CPU



of Threads

Get Performance on 32 Core Server



of Threads

P99 latency on 32 Core Server



micro-seconds

Experiment Setup

Server

- CPU @2.2GHz, 32 cores
- 20480KB cache & 144GB ram
- Key 20 bytes, value 100 bytes
- Prefix 12 bytes, 10 keys per prefix
- No compression
- 500M keys loaded with filluniquerandom (~64G data)
- Readwhilewriting with 10k/s write speed
- No backup performed

Take away

• By separating read-only and read-write data structures in the design, RocksDB makes performance scaling much easier

• Future work

- Iteration performance
- Write performance
- Improve cache efficiency for critical data structures